## Amendments to the Claims:

This listing of claims will replace all prior versions, and listings, of claims in the application.

## Listing of Claims:

(Currently Amended) An apparatus for allowing a user to model at least one
aspect <u>variation</u> of a software artifact by using extension types said-apparatus-comprising
a processor and a memory storing code accessible by the processor to provide extension
types, comprising:

a processor;

a memory accessible by the processor;

instructions contained in the memory and executable by the processor for allowing a user to obtain a controllable software artifact, said instructions comprising:

instructions for providing a user with a software artifact; and

instructions for allowing the user to add features to the software artifact;

wherein the features comprise extension types, each extension type comprising an ordered tuple of a plurality of element types, each of the element types corresponding to different class hierarchies; and

wherein said extension types are utilized to <u>implement simplify implementation of</u>
data classifications

- (Original) The apparatus according to Claim 1, wherein each extension type comprises an extension or variation of element types.
- (Original) The apparatus according to Claim 1, wherein said extension types are adapted to compose classes horizontally.
- (Original) The apparatus according to Claim 1, wherein each extension type is adapted to masquerade as any associated element type.
- (Original) The apparatus according to Claim 1, wherein each extension type is a subtype of its associated element types.
  - 6. (Original) The apparatus according to Claim 1, wherein:

each extension type has a size corresponding to the number of elements associated with the extension type; and

given two extension types  $\alpha$  and  $\beta$  , a sub-type relation  $\alpha < \beta$  is definable as follows:

$$|\alpha| >= |\beta|$$
; and

$$\alpha(0) <: \beta(0), \alpha(1) <: \beta(1), ... \alpha(|\beta|-1) <: \beta(|\beta|-1).$$

7. (Original) The apparatus according to Claim 1, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta < \alpha$ :

a method dispatch p.m comprises starting at the element type  $\beta(0)$  and walking up the class hierarchy of  $\beta(0)$  to find the closest m, wherein if m is not defined in the class hierarchy of  $\beta(0)$ , then m is sought in the  $\beta(1)$  class hierarchy and, if needed, in one or more iteratively successive class hierarchies, until found.

8. (Original) The apparatus according to Claim 1, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta <: \alpha$ :

a method dispatch  $p^*m$  comprises, for each element type  $\beta(i)$ , in the order  $i=0, ..., |\beta|-1$ , walking up the class hierarchy of  $\beta(i)$  to find the closest m in  $\beta(i)$  and dispatching the method m (if found), whereby a type error arises if m is not defined in at least one of the class hierarchies  $\beta(i)$ ,  $i=0, ..., |\beta|-1$ .

9. (**Original**) The apparatus according to Claim 1, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta <: \alpha$ :

a method dispatch p(1,3,4).m comprises reviewing only a class hierarchy of  $\beta(1)$ ,  $\beta(3)$ , and  $\beta(4)$ to find the closest m, wherein a type error arises if m is not defined in any of  $\beta(1)$ ,  $\beta(3)$ , or  $\beta(4)$ .

10. (Original) The apparatus according to Claim 1, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta <: \alpha$ :

a method dispatch p(1,3,4)\*m comprises reviewing only a class hierarchy of  $\beta(1)$ ,  $\beta(3)$ , and  $\beta(4)$ to find the closest m in  $\beta(i)$  and dispatching the method m if found, whereby a type error arises if in any of the class hierarchies to which  $\beta(1)$ ,  $\beta(3)$ , or  $\beta(4)$  belongs m is not defined.

11. (Currently Amended) A computer implemented method for allowing a user to model at least one aspect <u>variation</u> of a software artifact by using extension types, said method comprising;

providing the user with a controllable software artifact, said providing comprising:

providing the user with a software artifact; and

allowing the user to add features to the software artifact, wherein the features comprise the step of providing extension types, each extension type comprising an ordered tuple of a plurality of element types, each of the element types corresponding to different class hierarchies, wherein said extension types are stored in a memory of at least one general-purpose computer; and

wherein said extension types are utilized to <u>implement simplify implementation of</u>

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- 12. (Original) The method according to Claim 11, wherein each extension type comprises an extension or variation of element types.
- 13. (Original) The method according to Claim 11, wherein the extension types are adapted to compose classes horizontally.
- 14. (Original) The method according to Claim 11, wherein each extension type is adapted to masquerade as any associated element type.
- 15. (Original) The method according to Claim 11, wherein each extension type is a subtype of its associated element types.
  - 16. (Original) The method according to Claim 11, wherein:

each extension type has a size corresponding to the number of elements associated with the extension type; and

given two extension types  $\alpha$  and  $\beta$  , a sub-type relation  $\alpha < \beta$  is definable as follows:

$$|\alpha| >= |\beta|$$
; and

$$\alpha(0) <: \beta(0), \alpha(1) <: \beta(1), ... \alpha(|\beta|-1) <: \beta(|\beta|-1).$$

17. (Original) The method according to Claim 11, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta < \alpha$ :

a method dispatch p.m comprises starting at the element type  $\beta(0)$  and walking up the class hierarchy of  $\beta(0)$  to find the closest m, wherein if m is not defined in the class hierarchy of  $\beta(0)$ , then m is sought in the  $\beta(1)$  class hierarchy and, if needed, in one or more iteratively successive class hierarchies, until found.

18. (Original) The method according to Claim 11, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta$ <: $\alpha$ :

a method dispatch  $p^*m$  comprises, for each element type  $\beta(i)$ , in the order  $i=0, ..., |\beta|-1$ , walking up the class hierarchy of  $\beta(i)$  to find the closest m in  $\beta(i)$  and dispatching the method m (if found), whereby a type error arises if m is not defined in at least one of the class hierarchies  $\beta(i)$ ,  $i=0, ..., |\beta|-1$ .

19. (Original) The method according to Claim 11, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta <: \alpha$ :

a method dispatch p(1,3,4).m comprises reviewing only a class hierarchy of  $\beta(1)$ ,  $\beta(3)$ , and  $\beta(4)$ to find the closest m, wherein a type error arises if m is not defined in any of  $\beta(1)$ ,  $\beta(3)$ , or  $\beta(4)$ .

20. (Original) The method according to Claim 11, wherein, with  $\alpha$  being the extension type of a variable p and  $\beta$  being the runtime extension type of the object pointed by p, so that  $\beta <: \alpha$ :

a method dispatch p(1,3,4)\*m comprises reviewing only a class hierarchy of  $\beta(1)$ ,  $\beta(3)$ , and  $\beta(4)$  to find the closest m in  $\beta(i)$  and dispatching the method m if found, whereby a type error arises if in any of the class hierarchies to which  $\beta(1)$ ,  $\beta(3)$ , or  $\beta(4)$  belongs m is not defined.

21. (Currently Amended) A program storage device readable by machine, tangibly encoded with a program of instructions executable by a processor of the machine to perform method steps for allowing a user to model at least one variation of a software artifact by using extension types, said method steps comprising:

A data storage device readable by machine, comprising a data structure stored on the device, the data structure

providing the user with a controllable software artifact, said providing comprising:

providing the user with a software artifact; and

allowing the user to add features to the software artifact, wherein the features comprise being at least one extension type comprising an ordered tuple of a plurality of element types, each of the element types corresponding to different class hierarchies;

wherein said at least one extension type allows a user to model at least one aspect of a software artifact to simplify implementation of implement data classifications.